

MEMORY

OUR MARKETING SLIDE

- Kirk Pepperdine
 - Author of jPDM, a performance diagnostic model
 - Author of the original Java Performance Tuning workshop
- Co-founded Clarity
 - Building the smart generation of performance diagnostic tooling
 - Bring predictability into the diagnostic process
- Co-founded JCrete
 - The hottest unconference on the planet
- Java Champion(s)

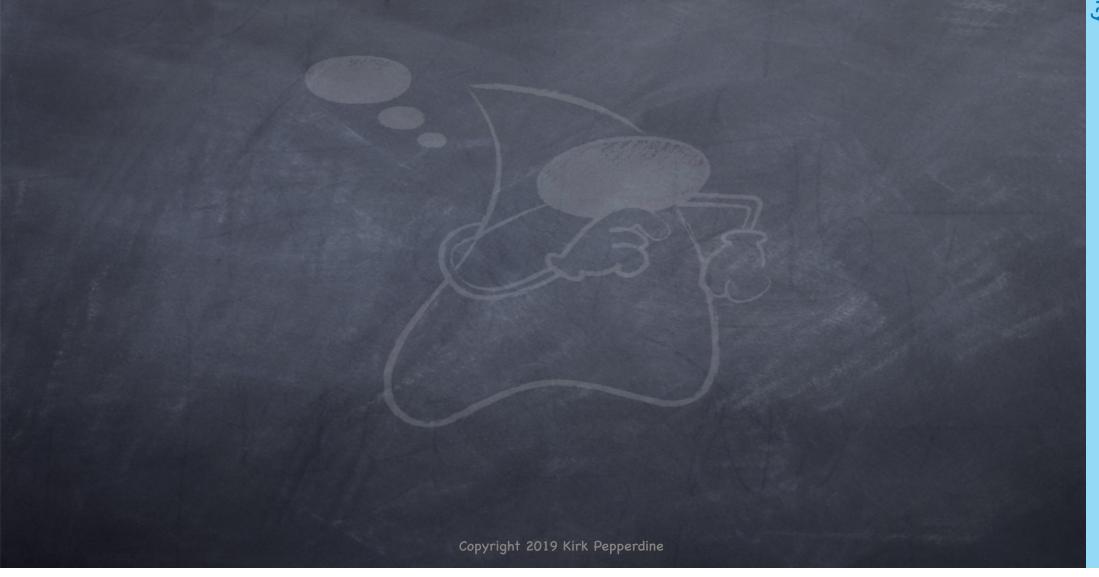
What is your performance trouble spot

INDUSTRY SURVEY



> 70% of all applications are bottlenecked on memory

and no,
Garbage Collection
is not a fault!!!!



Spring Boot

Cassandra

Cassandra or any big nosql solution

Apache Spark

Apache Spark or any big data framework



Log4J or any Java logging framework

JSON

JSON With almost any Marshalling protocol

ECom caching products

ECom caching products
Hibernate

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and so on

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then you are very likely in this 70%



WAR STORIES

- Reduced allocation rates from 1.8gb/sec to 0
 - > tps jumped from 400,000 to 25,000,000!!!
- Stripped all logging our of a transactional engine
 - Throughput jumped by a factor of 4x
- Wrapped 2 logging statements in a web socket framework
 - Memory churn reduced by a factor of 2

• and

ALLOCATION SITE

```
forms an allocation site

0: new #2 // class java/lang/Object
2: dup
4: invokespecial #1 // Method java/lang/Object."<init>":()V
```

- Allocation will (mostly) occur in Java heap
 - fast path
 - slow path
 - small objects maybe optimized to an on-stack allocation

JAVA HEAP



- Java Heap is made of;
 - Eden nursery
 - Survivor intermediate pool designed to delay promotion
 - ▶ Tenured to hold long lived data
- ▶ Each space contributes to a different set of problems
 - All affect GC overhead

EDEN ALLOCATIONS

top of heap pointer

```
top of heap pointer
```

```
Foo foo = new Foo();
Bar bar = new Bar();
byte[] array = new byte[N];
```

```
top of heap pointer

Foo foo = new Foo();
Bar bar = new Bar();
byte[] array = new byte[N];
```

```
Foo Bar

top of heap pointer

Foo foo = new Foo();

Bar bar = new Bar();

byte[] array = new byte[N];
```

Foo Bar byte[]

top of heap pointer

```
Foo foo = new Foo();
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```



top of heap pointer

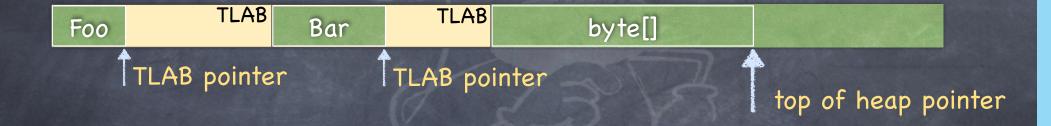
- In multi-threaded apps, top of heap pointer must be surrounded by barriers
 - single threads allocation
 - hot memory address
 - solved by stripping (Thread local allocation blocks)

TLAB ALLOCATION



- Assume 2 threads
 - each thread will have it's own (set of) TLAB(s)

TLAB ALLOCATIONS



- Thread 1 -> Foo foo = new Foo(); byte[] array = new byte[N];
 - byte[] doesn't fit in a TLAB
- Thread 2 -> Bar bar = new Bar();

TLAB WASTE %



- Threshold defining when to request a new TLAB
 - prevent buffer overflows
- waste up to 1% of a TLAB

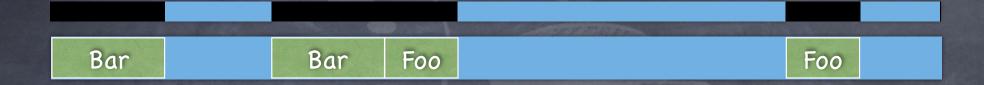
TLAB WASTE %



- Allocation failure to prevent buffer overflow
 - > some what expensive failure path

TENURED SPACE

Free List



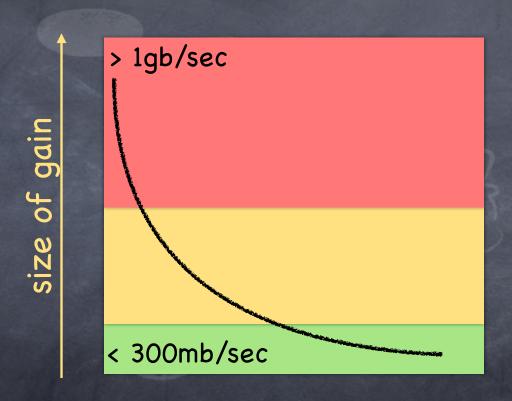
- Allocations in tenured make use of a free list
 - ▶ free list allocation is ~10x the cost of bump and run
- Data in tenured tends to be long lived
 - amount of data in tenured do affect GC pause times

- ▶ High memory churn rates
 - many temporary objects

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- Quickly fill Eden
 - frequent young gc cycles
 - speeds up aging
 - premature promotion
 - more frequent tenured cycles
 - increased copy costs
 - increased heap fragmentation
- Allocation is quick
 - quick * large number = slow

REDUCING ALLOCATIONS



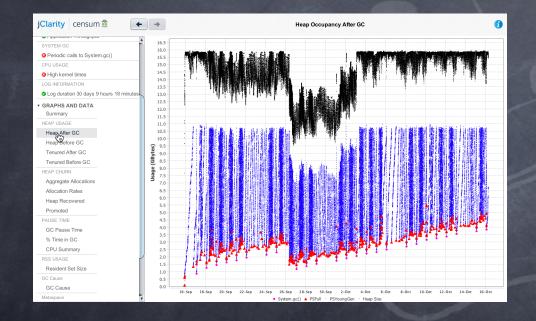
- ▶ High memory churn rates
 - many temporary objects
- Large live data set size
 - inflated live data set size
 - loitering

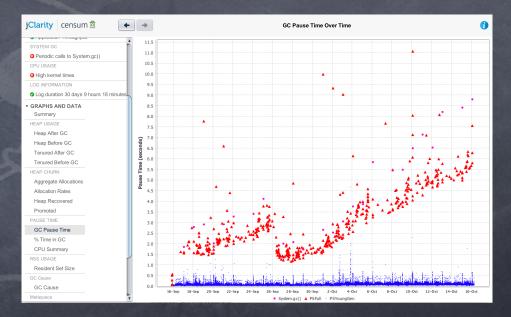
- ▶ High memory churn rates
 - many temporary objects

- Large live data set size
 - inflated live data set size
 - loitering

- inflated scan for root times
- reduced page locality
- Inflated compaction times
 - increase copy costs
 - likely less space to copy too

PAUSE VS OCCUPANCY





- ▶ High memory churn rates
 - many temporary objects

- Large live data set size
 - inflated live data set size
 - loitering
- Unstable live data set size
 - memory leak

- ▶ High memory churn rates
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- Eventually you run out of heap space
 - each app thread throws an OutOfMemoryError and terminates
 - JVM will shutdown with all nondaemon threads terminate

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Escape Analysis

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Demo time

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